Uninformed search strategies (Section 3.4)

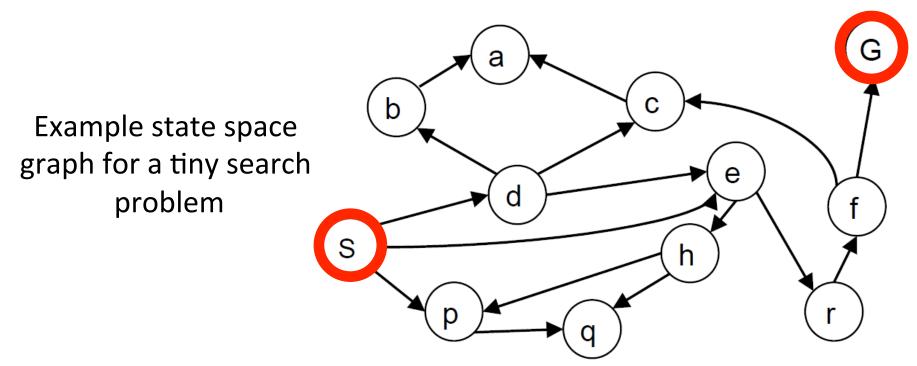


Uninformed search strategies

- A search strategy is defined by picking the order of node expansion
- Uninformed search strategies use only the information available in the problem definition
 - Breadth-first search
 - Depth-first search
 - Iterative deepening search
 - Uniform-cost search

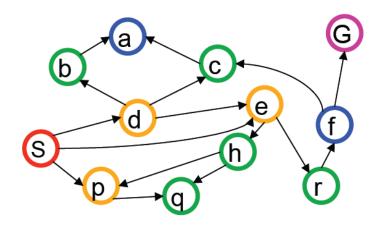
Breadth-first search

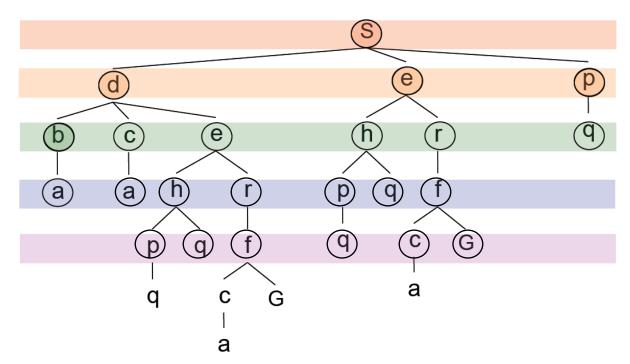
- Expand shallowest unexpanded node
- Implementation: frontier is a FIFO queue



Breadth-first search

Expansion order:
(S,d,e,p,b,c,e,h,r,q,a,a,h,r,p,q,f,p,q,f,q,c,G)



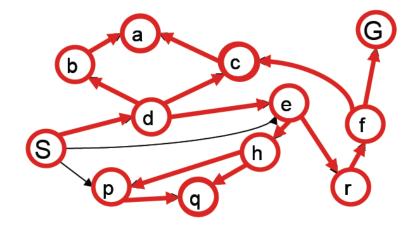


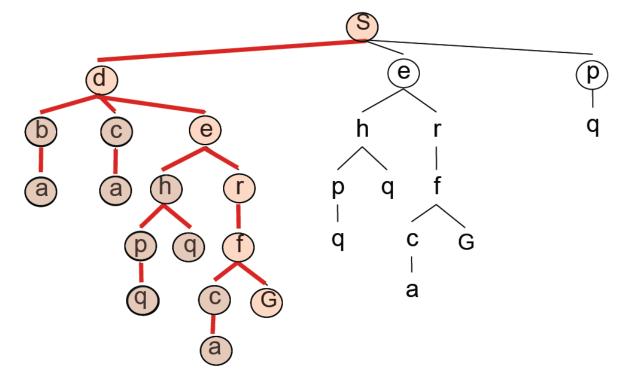
Depth-first search

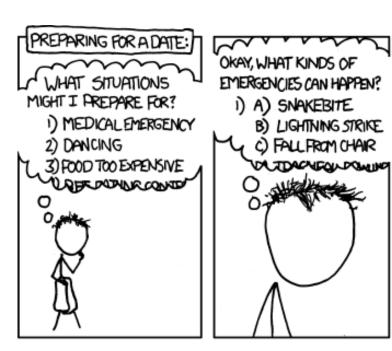
- Expand deepest unexpanded node
- Implementation: frontier is a LIFO queue

Depth-first search

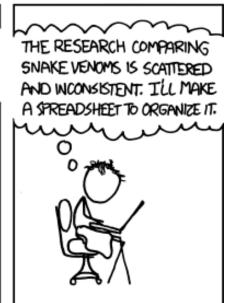
Expansion order:
(d,b,a,c,a,e,h,p,q,q,r,f,c,a,G)



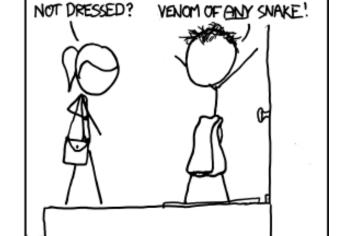








TAIPAN HAS THE DEADUIST

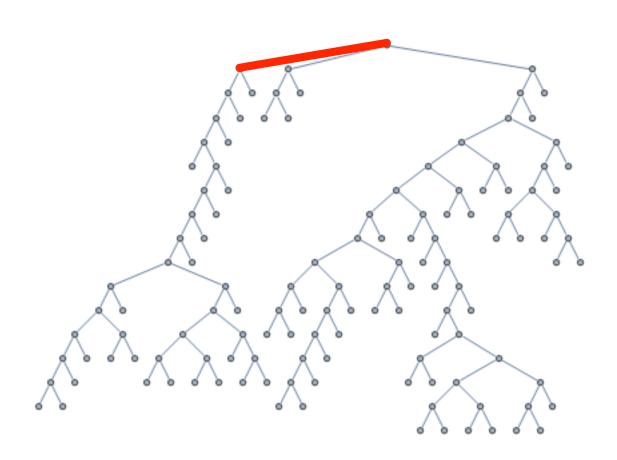


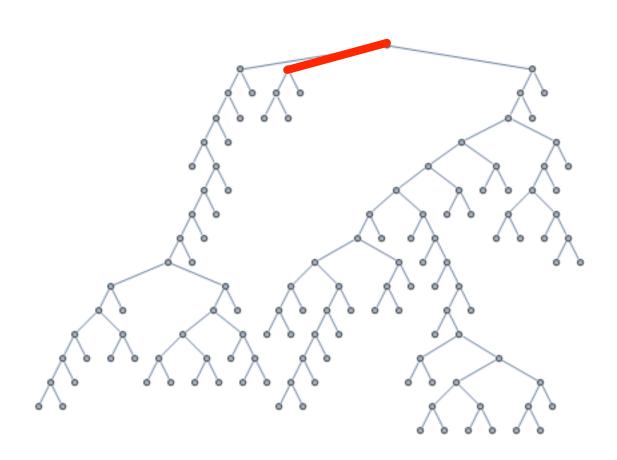
I'M HERE TO PICK BY LDSO, THE INLAND

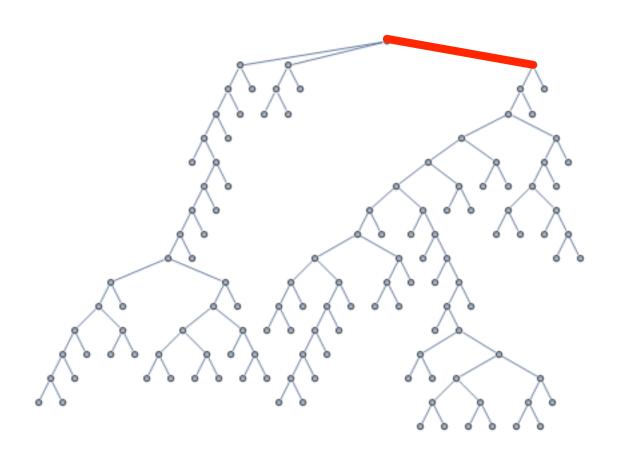
YOU'UP. YOU'RE

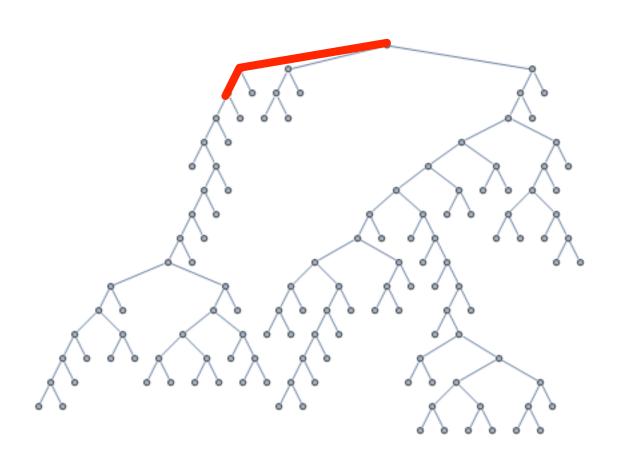
http://xkcd.com/761/

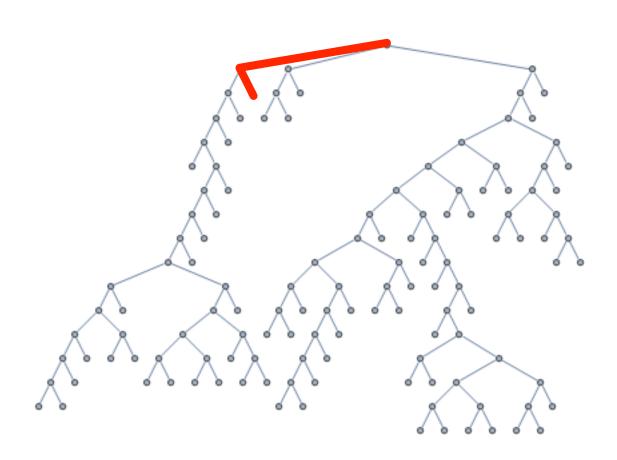
I REALLY NEED TO STOP USING DEPTH-FIRST SEARCHES.

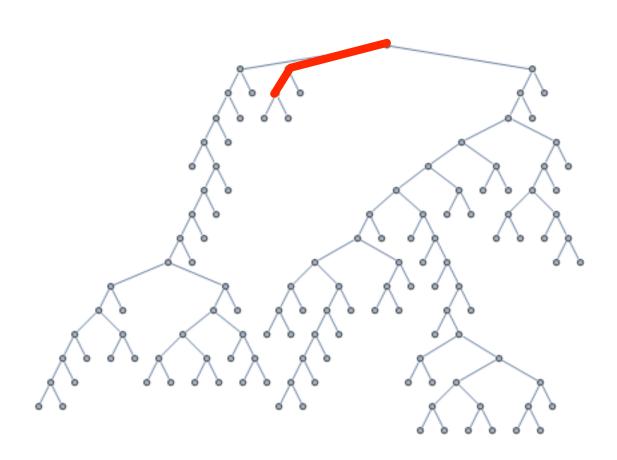


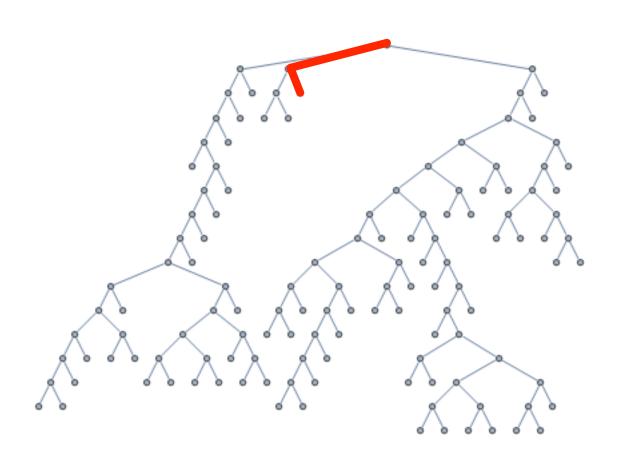


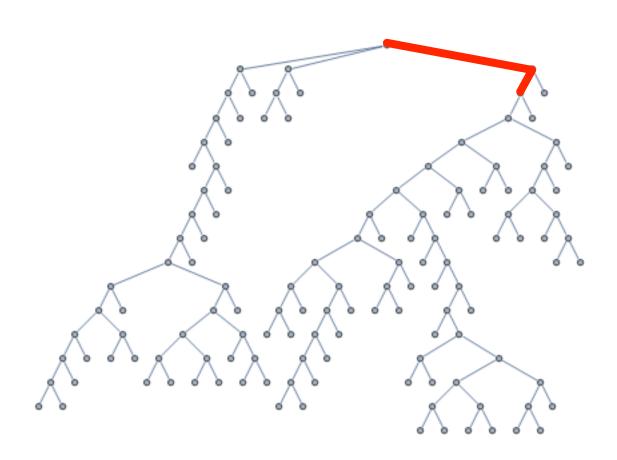


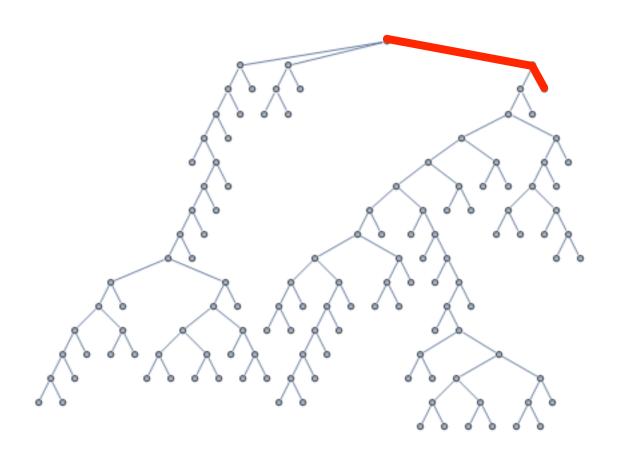


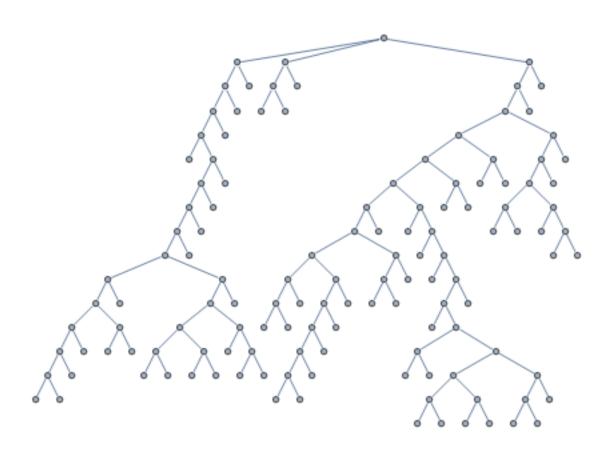


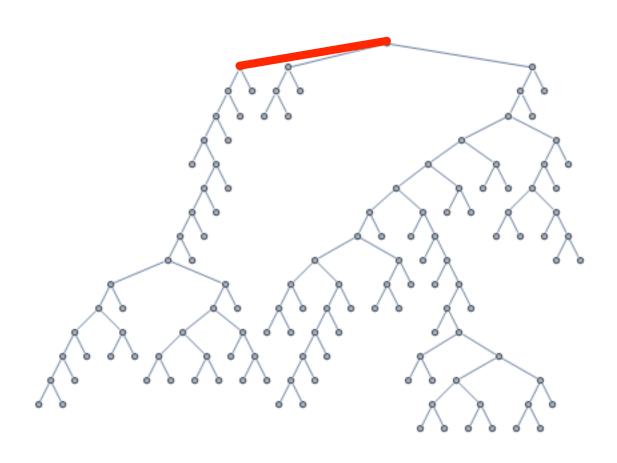


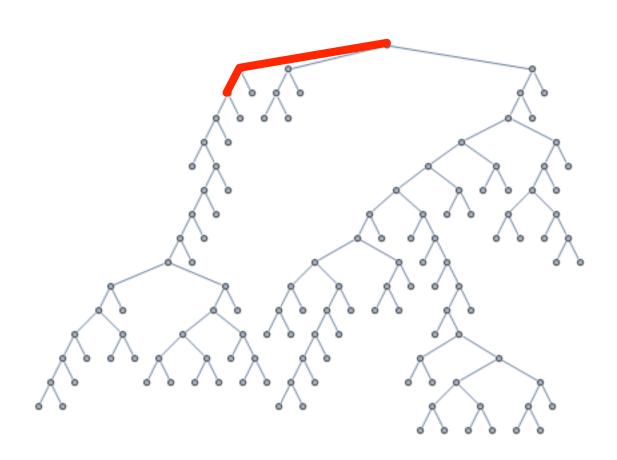


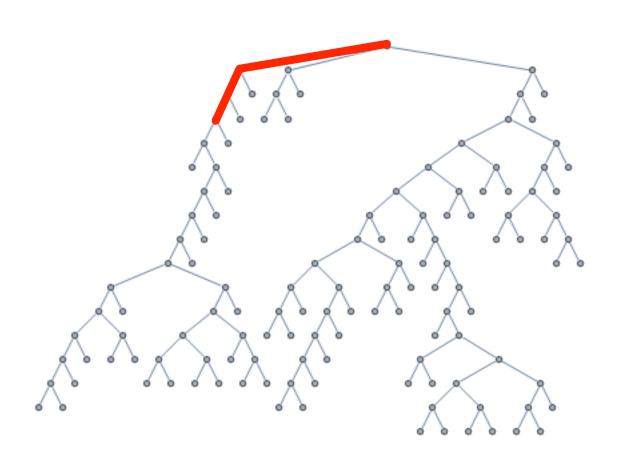


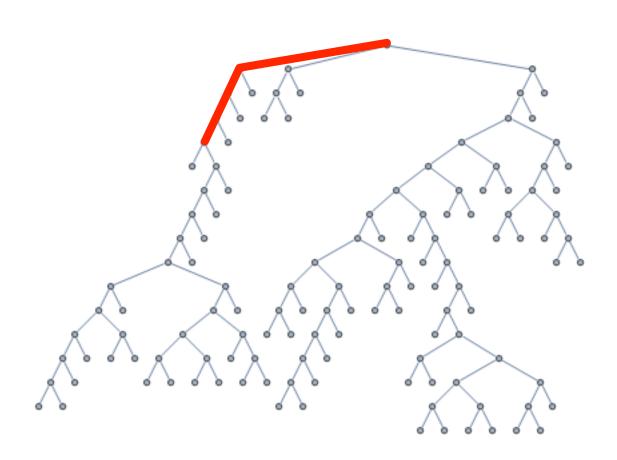


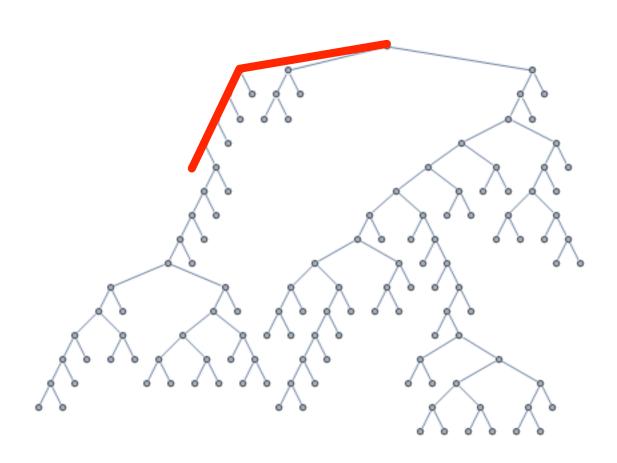


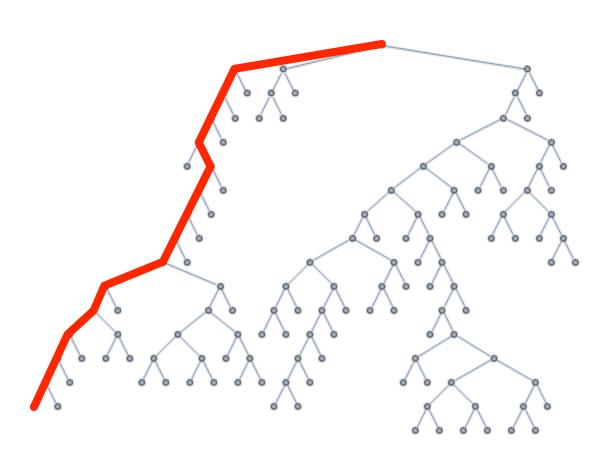


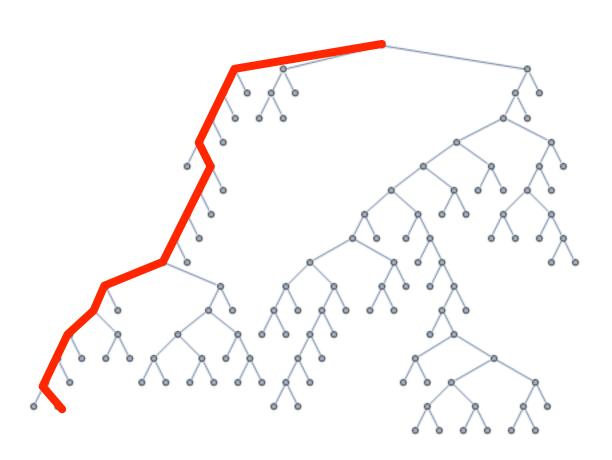


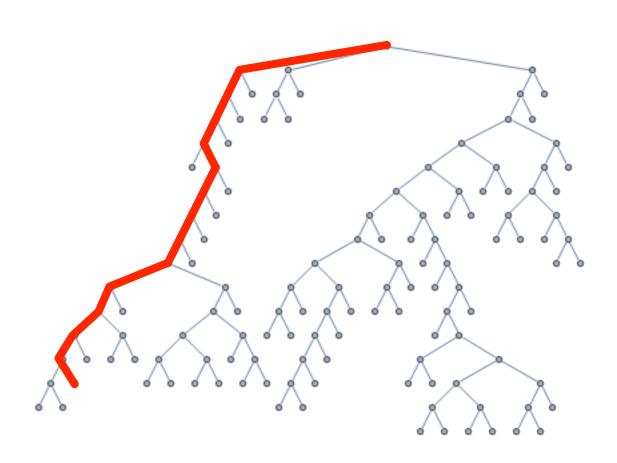


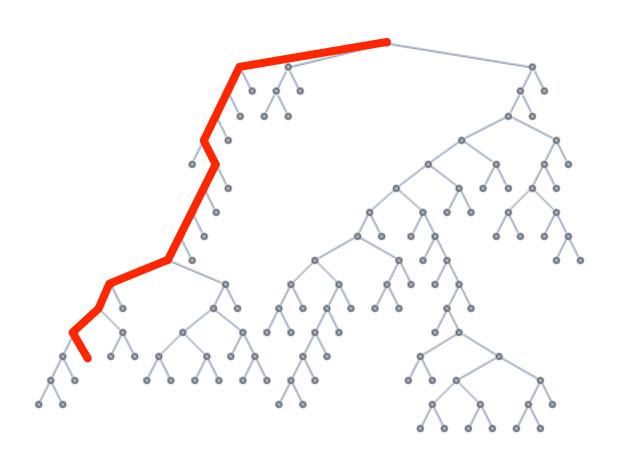


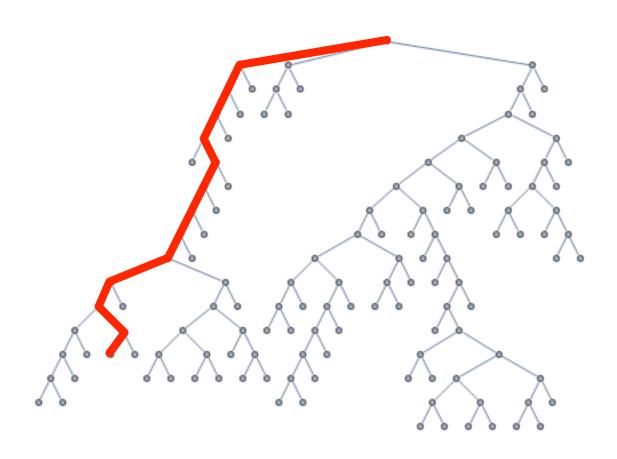


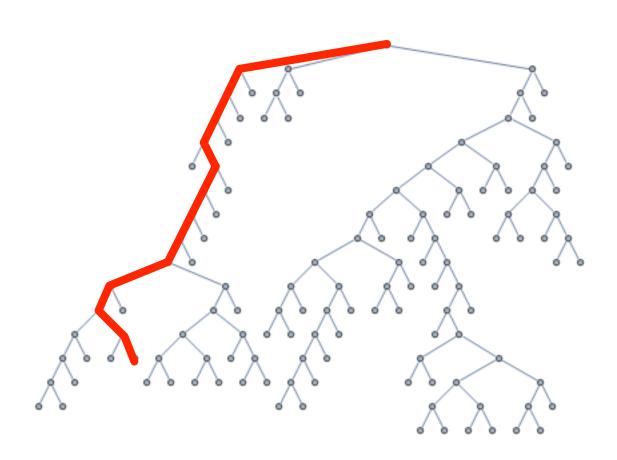


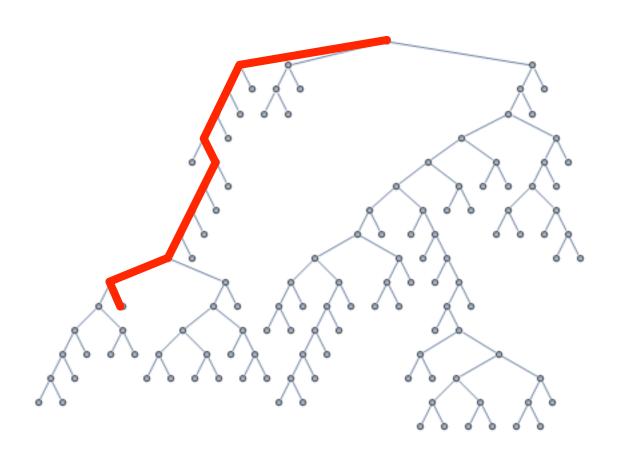












Analysis of search strategies

- Strategies are evaluated along the following criteria:
 - Completeness: does it always find a solution if one exists?
 - Optimality: does it always find a least-cost solution?
 - Time complexity: number of nodes generated
 - Space complexity: maximum number of nodes in memory
- Time and space complexity are measured in terms of
 - b: maximum branching factor of the search tree
 - d: depth of the optimal solution
 - m: maximum length of any path in the state space (may be infinite)

Properties of breadth-first search

Complete?

Yes (if branching factor **b** is finite)

Optimal?

Yes - if cost = 1 per step

Time?

Number of nodes in a b-ary tree of depth d: $O(b^d)$ (d is the depth of the optimal solution)

Space?

 $O(b^d)$

• Space is the bigger problem (more than time)

Properties of depth-first search

Complete?

Fails in infinite-depth spaces, spaces with loops Modify to avoid repeated states along path → complete in finite spaces

Optimal?

No – returns the first solution it finds

Time?

Could be the time to reach a solution at maximum depth m: $O(b^m)$ Terrible if m is much larger than d But if there are lots of solutions, may be much faster than BFS

Space?

O(bm), i.e., linear space!

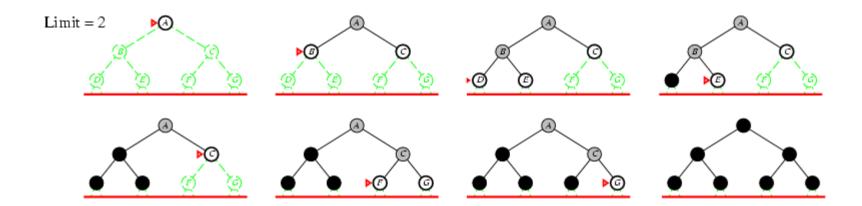
- Use DFS as a subroutine
 - 1. Check the root
 - 2. Do a DFS searching for a path of length 1
 - 3. If there is no path of length 1, do a DFS searching for a path of length 2
 - 4. If there is no path of length 2, do a DFS searching for a path of length 3...

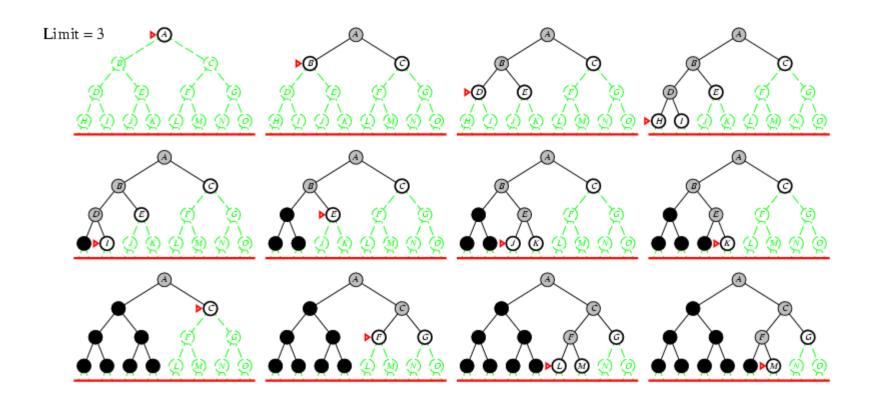
Limit = 0











Properties of iterative deepening search

Complete?

Yes

Optimal?

Yes, if step cost = 1

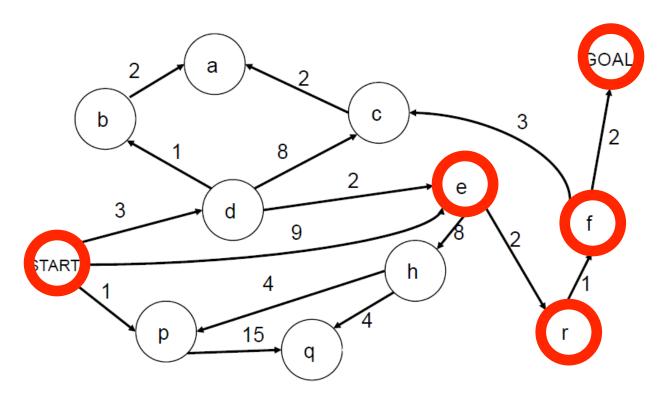
Time?

$$(d+1)b^0 + db^1 + (d-1)b^2 + ... + b^d$$

Space?

O(bd)

Search with varying step costs



 BFS finds the path with the fewest steps, but does not always find the cheapest path

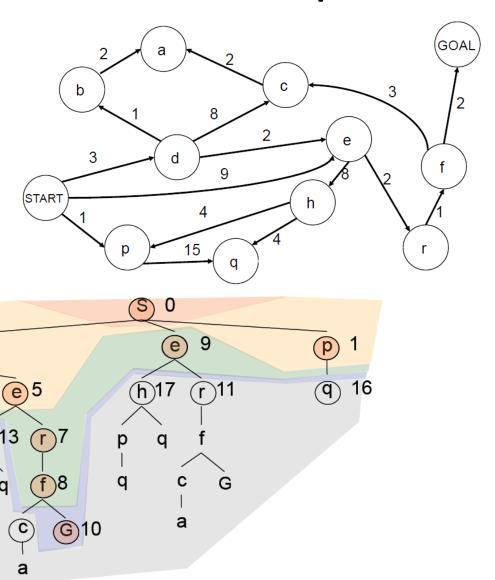
Uniform-cost search

- For each frontier node, save the total cost of the path from the initial state to that node
- Expand the frontier node with the lowest path cost
- Implementation: frontier is a priority queue ordered by path cost
- Equivalent to BFS if step costs all equal
- Equivalent to Dijkstra's algorithm in general

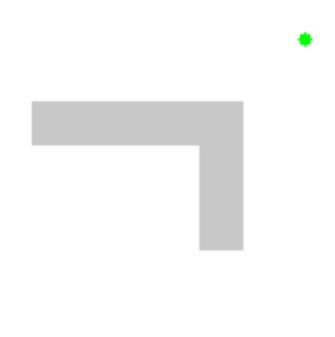
Uniform-cost search example

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Expansion order: (S,p,d,b,e,a,r,f,e,G)



Another example of uniform-cost search



Source: Wikipedia

Properties of uniform-cost search

Complete?

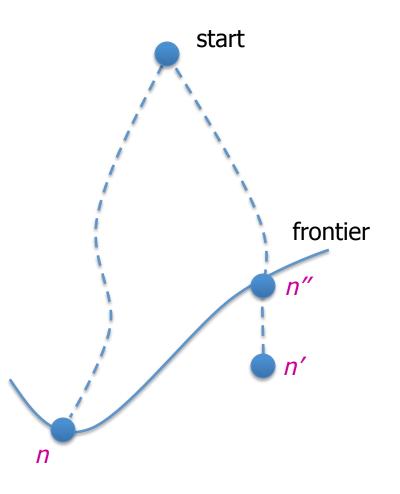
Yes, if step cost is greater than some positive constant ε (we don't want infinite sequences of steps that have a finite total cost)

Optimal?

Yes

Optimality of uniform-cost search

- Graph separation property: every path from the initial state to an unexplored state has to pass through a state on the frontier
 - Proved inductively
- Optimality of UCS: proof by contradiction
 - Suppose UCS terminates at goal state n with path cost g(n) but there exists another goal state n' with g(n') < g(n)
 - By the graph separation property, there must exist a node n'' on the frontier that is on the optimal path to n'
 - But because g(n'') ≤ g(n') < g(n), n'' should have been expanded first!



Properties of uniform-cost search

Complete?

Yes, if step cost is greater than some positive constant ε (we don't want infinite sequences of steps that have a finite total cost)

Optimal?

Yes – nodes expanded in increasing order of path cost

Time?

Number of nodes with path cost \leq cost of optimal solution (C^*), $O(b^{C^*/\varepsilon})$

This can be greater than $O(b^d)$: the search can explore long paths consisting of small steps before exploring shorter paths consisting of larger steps

Space?

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O(b^{C^*/\varepsilon})
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Review: Uninformed search strategies

Algorithm	Complete?	Optimal?	Time complexity	Space complexity
BFS	Yes	If all step costs are equal	O(b ^d)	O(b ^d)
DFS	No	No	O(b ^m)	O(bm)
IDS	Yes	If all step costs are equal	O(b ^d)	O(bd)
UCS	Yes	Yes	Number of node	es with g(n) ≤ C*

b: maximum branching factor of the search tree

d: depth of the optimal solution

m: maximum length of any path in the state space

C*: cost of optimal solution

g(n): cost of path from start state to node n